On Kronecker Products, Tensor Products and Matrix Differential Calculus

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ON KRONECKER PRODUCTS, TENSOR PRODUCTS
AND MATRIX DIFFERENTIAL CALCULUS

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The algebra of the Kronecker products of matrices is recapitulated using a
notation that reveals the tensor structures of the matrices. It is claimed that
many of the difficulties that are encountered in working with the algebra
can be alleviated by paying close attention to the indices that are concealed
beneath the conventional matrix notation.

The vectorisation operations and the commutation transformations
that are common in multivariate statistical analysis alter the positional
relationship of the matrix elements. These elements correspond to numbers
that are liable to be stored in contiguous memory cells of a computer, which
should remain undisturbed.

It is suggested that, in the absence of an adequate index notation that
enables the manipulations to be performed without disturbing the data,
even the most clear-headed of computer programmers is liable to perform
wholly unnecessary and time-wasting operations that shift data between
memory cells.

1. Introduction

One of the bugbears of multivariate statistical analysis is the need to differen-
tiate complicated likelihood functions in respect of their matrix arguments. To
achieve this, one must resort to the theory of matrix differential calculus, which
entails the use of Kronecker products, vectorisation operators and commutation
matrices.

A brief account of the requisite results was provided by Pollock (1979),
who described a theory that employs vectorised matrices. Relationships were
established with the classical theory of Dwyer and McPhail (1948) and Dwyer
(1967), which dealt with scalar functions of matrix augments and matrix func-
tions of scalar arguments. A contemporaneous account was given by Henderson
and Searle (1979), and similar treatments were provided, at roughly the same
time, by Balestra (1976) and by Rogers (1980).

A more extensive account of matrix differential calculus, which relies ex-
clusively on vectorised matrices, was provided by the text of Magnus and
Neudecker (1988). This has become a standard reference. More recent ac-
counts of matrix differential calculus have been provided by Turkington (2002)
and by Harville (2008).

Notwithstanding this ample provision of sources, there continues to be
confusion and difficulty in this area. A recent paper of Magnus (2010), which
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testifies to these problems, cites several instances in which inappropriate definitions of matrix derivatives have been adopted, and it gives examples that highlight the perverse effects of adopting such definitions.

One is bound to wonder why the difficulties persist. An answer to this conundrum, which is offered in the present paper, points to the absence of an accessible account that reveals the tensorial structures that underlie the essential results of matrix differential calculus and of the associated algebra.

Part of the problem lies in the fact that the conventional matrix notation, which has been employed, almost exclusively, to convey the theory of matrix differential calculus, conceals the very things that need to be manipulated, which are matrix elements bearing strings of indices that are ordered according to some permutation of a lexicographic ordering. To overcome the problem, we shall adopt a notation that reveals the indices.

This paper aims to avoid unnecessary abstractions by dealing only with concrete objects, which are matrices and their elements. A more abstract approach would employ the concepts and the notations of vector spaces and of their tensor products. Such an approach to the theory is exemplified by the texts in multilinear algebra of Bourbaki (1958), Greub (1967) and Marcus (1973). A primary source of multilinear algebra and differential geometry is Cartan (1952). One advantage of the approach that we shall take is that it is well adapted to the requirements of computer programming.

2. Bases for Vector Spaces

Consider an identity matrix of order $N$, which can be written as follows:

\[
\begin{bmatrix}
1 & 0 & \cdots & 0 \\
0 & 1 & \cdots & 0 \\
\vdots & \vdots & \ddots & \vdots \\
0 & 0 & \cdots & 1
\end{bmatrix}
= \begin{bmatrix}
e_1 \\
e_2 \\
\vdots \\
e_N
\end{bmatrix}
\]

On the LHS, the matrix is expressed as a collection of column vectors, denoted by $e_i; i = 1, 2, \ldots, N$, which form the basis of an ordinary $N$-dimensional Euclidean space, which is the primal space. On the RHS, the matrix is expressed as a collection of row vectors $e^j; j = 1, 2, \ldots, N$, which form the basis of the conjugate dual space. In general, vectors bearing a single superscript, including a prime, are to be regarded as row vectors.

The basis vectors can be used in specifying arbitrary vectors in both spaces. In the primal space, there is the column vector

\[
a = \sum_i a_i e_i = (a_i e_i),
\]

and in the dual space, there is the row vector

\[
b' = \sum_j b_j e^j = (b_j e^j).
\]
Here, on the RHS, there is a notation that replaces the summation signs by parentheses. When a basis vector is enclosed by parentheses, summations are to be taken in respect of the index or indices that it carries. Usually, such indices will be associated with scalar elements that will also be found within the parentheses. The advantage of this notation will become apparent at a later stage, when the summations are over several indices.

A vector in the primary space can be converted to a vector in the conjugate dual space, and vice versa, by the operation of transposition. Thus, \( a' = (a_i e^i) \) is formed from \( a = (a_i e_i) \) via the conversion \( e_i \to e^i \), whereas \( b = (b_j e_j) \) is formed from \( b' = (b_j e^j) \) via the conversion \( e^j \to e_j \).

### 3. Elementary Tensor Products

A tensor product of two vectors is an outer product that entails the pairwise products of the elements of both vectors. Consider two primal vectors

\[
\begin{align*}
 a &= [a_t; t = 1, \ldots, T] = [a_1, a_2, \ldots, b_T]' \quad \text{and} \\
 b &= [b_j; j = 1, \ldots, M] = [b_1, b_2, \ldots, b_M]',
\end{align*}
\]

which need not be of the same order. Then, two kinds of tensor products can be defined. First, there are covariant tensor products. The covariant product of \( a \) and \( b \) is a column vector in a primal space:

\[
(5) \quad a \otimes b = \sum_t \sum_j a_t b_j (e_t \otimes e_j) = (a_t b_j e_{tj}).
\]

Here, the elements are arrayed in a long column in an order that is determined by the lexicographic variation of the indices \( t \) and \( j \). Thus, the index \( j \) undergoes a complete cycle from \( j = 1 \) to \( j = M \) with each increment of the index \( t \) in the manner that is familiar from dictionary classifications. Thus

\[
(6) \quad a \otimes b = \begin{bmatrix} a_1 b_1 \\ a_2 b_2 \\ \vdots \\ a_T b_T \end{bmatrix} = [a_1 b_1, \ldots, a_1 b_M, a_2 b_1, \ldots, a_2 b_M, \ldots, a_T b_1, \ldots, a_T b_M]'.
\]

A covariant tensor product can also be formed from the row vectors \( a' \) and \( b' \) of the dual space. Thus, there is

\[
(7) \quad a' \otimes b' = \sum_t \sum_j a_t b_j (e^t \otimes e^j) = (a_t b_j e^{tj}).
\]

It will be observed that this is just the transpose of \( a \otimes b \). That is to say,

\[
(8) \quad (a \otimes b)' = a' \otimes b' \quad \text{or, equivalently,} \quad (a_t b_j e_{tj})' = (a_t b_j e^{tj}).
\]
The order of the vectors in a covariant tensor product is crucial, since, as one can easily verify, it is the case that

\begin{equation}
    a \otimes b \neq b \otimes a \quad \text{and} \quad a' \otimes b' \neq b' \otimes a'.
\end{equation}

The second kind of tensor product of the two vectors is a so-called contravariant tensor product:

\begin{equation}
    a \otimes b' = b' \otimes a = \sum_{t} \sum_{j} a_{t}b_{j}(e_{t} \otimes e^{j}) = (a_{t}b_{j}e^{j}).
\end{equation}

This is just the familiar matrix product $ab'$, which can be written variously as

\begin{equation}
    \begin{bmatrix}
        a_{1} b' \\
        a_{2} b' \\
        \vdots \\
        a_{T} b'
    \end{bmatrix}
    = [b_{1} a \ b_{2} a \ \cdots \ b_{M} a]
    = \begin{bmatrix}
        a_{1} b_{1} & a_{1} b_{2} & \cdots & a_{1} b_{M} \\
        a_{2} b_{1} & a_{2} b_{2} & \cdots & a_{2} b_{M} \\
        \vdots & \vdots & \ddots & \vdots \\
        a_{T} b_{1} & a_{T} b_{2} & \cdots & a_{T} b_{M}
    \end{bmatrix}.
\end{equation}

According to (10), the ordering of the vectors within such a binary contravariant tensor product is immaterial, albeit that $ab' \neq b'a$.

Observe that

\begin{equation}
    (a \otimes b')' = a' \otimes b \quad \text{or, equivalently,} \quad (a_{t}b_{j}e^{j})' = (a_{t}b_{j}e^{j}).
\end{equation}

We now propose to dispense with the summation signs and to write the various vectors as follows:

\begin{equation}
    a = (a_{t}e_{t}), \quad a' = (a_{t}e^{t}) \quad \text{and} \quad b = (a_{j}e_{j}), \quad b' = (b_{j}e^{j}).
\end{equation}

As before, the convention here is that, when the products are surrounded by parentheses, summations are to be taken in respect of the indices that are associated with the basis vectors.

The convention can be applied to provide summary representations of the products under (5), (7) and (10):

\begin{align*}
    a \otimes b &= (a_{t}e_{t}) \otimes (b_{j}e_{j}) = (a_{t}b_{j}e_{tj}), \\
    a' \otimes b' &= (a_{t}e^{t}) \otimes (b_{j}e^{j}) = (a_{t}b_{j}e^{tj}), \\
    a \otimes b' &= (a_{t}e_{t}) \otimes (b_{j}e^{j}) = (a_{t}b_{j}e^{j}).
\end{align*}

Such products are described as decomposable tensors.

4. Non-decomposable Tensor Products

Non-decomposable tensors are the result of taking weighted sums of decomposable tensors. Consider an arbitrary matrix $X = [x_{tj}]$ of order $T \times M$. This can
be expressed as the following weighted sum of the contravariant tensor products formed from the basis vectors:

\[ X = (x_{tj} e_t^j) = \sum_t \sum_j x_{tj} (e_t \otimes e^j).\]  

(17)

The indecomposability lies in the fact that the elements \( x_{tj} \) cannot be written as the products of an element indexed by \( t \) and an element indexed by \( j \).

From \( X = (x_{tj} e_t^j) \), the following associated tensors products may be derived:

\[ X' = (x_{tj} e_t^j), \]  

(18)

\[ X^r = (x_{tj} e^{tj}), \]  

(19)

\[ X^c = (x_{tj} e_{jt}). \]  

(20)

Here, \( X' \) is the transposed matrix, whereas \( X^c \) is a long column vector and \( X^r \) is a long row vector. Notice that, in forming \( X^c \) and \( X^r \) from \( X \), the index that moves assumes a position at the head of the string of indices to which it is joined. It will be observed that whereas the indices of the elements of \( X^r \) follow a lexicographic ordering, with the leading index as the principal classifier, those of \( X^c \) follow a reverse lexicographic ordering.

The superscript letters \( c \) and \( r \) denote a pair of so-called vectorisation operators. It has become conventional to denote \( X^c \) by \( \text{vec} X \). Turkington (2000 and 2002) has used the notation \( \text{devec} X = (\text{vec} X^r)' \) to denote \( X^r \) and to indicate its relationship with \( \text{vec} X \).

It is evident that

\[ X^r = X'^{cr} \quad \text{and} \quad X^c = X'^{rc}. \]  

(21)

Thus, it can be seen that \( X^c \) and \( X^r \) are not related to each other by simple transpositions.

The transformation that effects the reversal of the ordering of the two indices, and which thereby reverses the order of the vectors in a covariant tensor product, was described by Pollock (1979) as the tensor commutator. The transformation was denoted by a capital \( T \) inscribed in a circle. The copyright symbol is also an appropriate choice of notation, which leads one to write

\[ X'^r = X'^c = \odot X^c \quad \text{and} \quad X^c = \odot X'^c. \]  

(22)

It will be shown in Section 7 that the transformation \( \odot \) corresponds to an orthonormal matrix. This was described by Magnus and Neudecker (1979) as the commutation matrix, which they denoted by \( K \).

**Example.** Consider the equation

\[ y_{tj} = \mu + \gamma_t + \delta_j + \varepsilon_{tj}. \]  

(23)
wherein \( t = 1, \ldots, T \) and \( j = 1, \ldots, M \). This relates to a two-way analysis of variance. For a concrete interpretation, we may imagine that \( y_{tj} \) is an observation taken at time \( t \) in the \( j \)th region. Then, the parameter \( \gamma_t \) represents an effect that is common to all observations taken at time \( t \), whereas the parameter \( \delta_j \) represents a characteristic of the \( j \)th region that prevails through time.

In ordinary matrix notation, the set of TM equations becomes

\[
Y = \mu T \nu^*_M + \gamma T \nu^*_M + \nu T \delta^* + \mathcal{E},
\]

where \( Y = [y_{tj}] \) and \( \mathcal{E} = [\varepsilon_{tj}] \) are matrices of order \( T \times M \), \( \gamma = [\gamma_1, \ldots, \gamma_T]^T \) and \( \delta = [\delta_1, \ldots, \delta_M]^T \) are vectors of orders \( T \) and \( M \) respectively, and \( \nu T \) and \( \nu M \) are vectors of units whose orders are indicated by their subscripts. In terms of the index notation, the TM equations are represented by

\[
(y_{tj} e^*_i) = \mu (e^*_i) + (\gamma_t e^*_i) + (\delta_j e^*_i) + (\varepsilon_{tj} e^*_i).
\]

An illustration is provided by the case where \( T = M = 3 \). Then equations (24) and (25) represent the following structure:

\[
\begin{bmatrix}
y_{11} & y_{12} & y_{13} \\
y_{21} & y_{22} & y_{23} \\
y_{31} & y_{32} & y_{33}
\end{bmatrix}
= \mu \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}
+ \begin{bmatrix} \gamma_1 & \gamma_1 & \gamma_1 \\ \gamma_2 & \gamma_2 & \gamma_2 \\ \gamma_3 & \gamma_3 & \gamma_3 \end{bmatrix}
+ \begin{bmatrix} \delta_1 & \delta_2 & \delta_3 \\ \delta_1 & \delta_2 & \delta_3 \\ \delta_1 & \delta_2 & \delta_3 \end{bmatrix}
+ \begin{bmatrix} \varepsilon_{11} & \varepsilon_{12} & \varepsilon_{13} \\ \varepsilon_{21} & \varepsilon_{22} & \varepsilon_{23} \\ \varepsilon_{31} & \varepsilon_{32} & \varepsilon_{33} \end{bmatrix}.
\]

5. Multiple Tensor Products

The tensor product entails an associative operation that combines matrices or vectors of any order. Let \( B = [b_{ij}] \) and \( A = [a_{ki}] \) be arbitrary matrices of orders \( t \times n \) and \( s \times m \) respectively. Then, their tensor product \( B \otimes A \), which is also known as a Kronecker product, is defined in terms of the index notation by writing

\[
(b_{ij} e^*_k) \otimes (a_{ki} e^*_l) = (b_{ij} a_{ki} e^{jl}_{lk}).
\]

Here, \( e^{ji}_{lk} \) stands for a matrix of order \( st \times mn \) with a unit in the row indexed by \( lk \)—the \( ((l-1)s+k) \)th row—and in the column indexed by \( ji \)—the \( ((j-1)m+l) \)th column—and with zeros elsewhere.

In the matrix array, the row indices \( lk \) follow a lexicographic order, as do the column indices \( ji \). Also, the indices \( lk \) are not ordered relative to the indices \( ji \). That is to say,

\[
e^{ji}_{lk} = e^i \otimes e^j \otimes e^k \otimes e^l = e^i \otimes e^j \otimes e^l \otimes e^k = e^j \otimes e^i \otimes e^k \otimes e^l = e^i \otimes e^j \otimes e^k \otimes e^l = e^j \otimes e^i \otimes e^l \otimes e^k.
\]
The virtue of the index notation is that it makes no distinction amongst these various products on the RHS—unless a distinction can be found between such expressions as $e^{ij}_{ik}$ and $e^{ij}_{ik}$.

For an example, consider the Kronecker of two matrices as follows:

$$
\begin{bmatrix}
  b_{11} & b_{12} \\
  b_{21} & b_{22}
\end{bmatrix} \otimes \begin{bmatrix}
  a_{11} & a_{12} \\
  a_{21} & a_{22}
\end{bmatrix} = \begin{bmatrix}
  b_{11} & a_{11} & a_{12} & b_{12} & a_{11} & a_{12} \\
  b_{21} & a_{21} & a_{22} & b_{22} & a_{21} & a_{22}
\end{bmatrix}
$$

(29)

Here, it can be seen that the composite row indices $lk$, associated with the elements $b_{ij}a_{ki}$, follow the lexicographic sequence \{11, 12, 21, 22\}. The column indices follow the same sequence.

6. Compositions

In order to demonstrate the rules of matrix composition, let us consider the matrix equation

$$
Y = AXB',
$$

(30)

which can be construed as a mapping from $X$ to $Y$. In the index notation, this is written as

$$
(y_{kit}e^l_k) = (a_{kit}e^l_k)(x_{ij}e^i_j)(b_{lj}e^l_j)
$$

(31)

$$
= (\{a_{kit}x_{ij}b_{lj}\}e^l_k).
$$

Here, there is

$$
\{a_{kit}x_{ij}b_{lj}\} = \sum_i \sum_j a_{kit}x_{ij}b_{lj};
$$

(32)

which is to say that the braces surrounding the expression on the LHS are to indicate that summations are taken with respect to the repeated indices $i$ and $j$, which are associated with the basis vectors. The operation of composing two factors depends upon the cancellation of a superscript (column) index, or string of indices, in the leading factor with an equivalent subscript (row) index, or string of indices, in the following factor.

The matrix equation of (30) can be vectorised in a variety of ways. In order to represent the mapping from $X^c = (x_{ij}e_{ji})$ to $Y^c = (y_{kit}e_{lk})$, we may write

$$
(y_{kit}e_{lk}) = (\{a_{kit}x_{ij}b_{lj}\}e_{lk})
$$

(33)

$$
= (a_{kit}b_{lj}e^i_{lk})(x_{ij}e^j_{ji}).
$$

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Notice that the product $a_{ki}b_{ij}$ within $(a_{ki}b_{ij}e_{ik}^{ji})$ does not need to be surrounded by braces, since it contains no repeated indices. Nevertheless, there would be no harm in writing $\{a_{ki}b_{ij}\}$.

The matrix $(a_{ki}b_{ij}e_{ik}^{ji})$ is decomposable. That is to say

$$ (a_{ki}b_{ij}e_{ik}^{ji}) = (b_{ij}e_{ij}^c) \otimes (a_{ki}e_{ik}^c) = B \otimes A; $$

and, therefore, the vectorised form of equation (30) is

$$ Y^c = (AXB')^c = (B \otimes A)X^c. $$

Other allied results that readily derived from (35) are

$$ Y'^c = Y'^T = (A \otimes B)X'^c, $$

$$ Y'^c = Y'^T = X'^c(B' \otimes A'), $$

$$ Y'^c = X'^c(A' \otimes B'). $$

Another result, which may be noted in this context, concerns the trace of a matrix product. If $A = (a_{ij}e_{ij}^c)$ and $B = (b_{ij}e_{ij}^c)$, then

$$ \text{Trace}(AB) = \{a_{ij}b_{ji}\} = (a_{ij}e_{ij}^c)(b_{ji}e_{ij}^c) = A'B^c. $$

Here, we are invoking the convention that $\{a_{ij}b_{ji}\}$ denotes the sum over the repeated indices $i$ and $j$.

**Example.** The equation under (25), which relates to a two-way analysis of variance, can be vectorised to give

$$ (y_{i}e_{j}) = \mu(e_{j}) + (e_{j1}^t)(\gamma_{i}e_{i}) + (e_{ij}^t)(\delta_{j}e_{j}) + (\epsilon_{ij}e_{jt}). $$

Using the notation of the Kronecker product, this can also be rendered as

$$ Y^c = \mu(\iota_{M} \otimes \nu_{T}) + (\iota_{M} \otimes I_{T})\gamma + (I_{M} \otimes \nu_{T})\delta + \Sigma^c. $$

The latter can also be obtained by applying the rule of (35) to equation (24). The various elements of (24) have been vectorised as follows:

$$ (\mu\iota_{M}^{T})^c = (\nu_{T}\mu_{M}^{T})^c = (\iota_{M} \otimes \nu_{T})\mu, $$

$$ (\gamma\iota_{M}^{T})^c = (I_{T}\gamma_{M}^{T})^c = (\iota_{M} \otimes I_{T})\gamma, $$

$$ (\nu_{T}\delta)^c = (\nu_{T}\delta^{T}I_{M})^c = (I_{M} \otimes \nu_{T})\delta^c, \quad \delta^c = \delta. $$

Also, there is $(\iota_{M} \otimes \nu_{T})\mu = \mu(\iota_{M} \otimes \nu_{T})$, since $\mu$ is a scalar element that can be transposed or freely associated with any factor of the expression.
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In comparing (40) and (41), we see, for example, that \((e^t_j) = (e_j) \otimes (e^t_j) = \iota_M \otimes \iota_T\). We recognise that \((e^t_j)\) is the sum over the index \(t\) of the matrices of order \(T\) which have a unit in the \(t\)th diagonal position and zeros elsewhere; and this sum amounts, of course, to the identity matrix of order \(T\).

The vectorised form of equation (26) is

\[
\begin{bmatrix}
  y_{11} \\
y_{21} \\
y_{31} \\
y_{12} \\
y_{22} \\
y_{32} \\
y_{13} \\
y_{23} \\
y_{33}
\end{bmatrix}
= \begin{bmatrix}
  1 & 0 & 0 & 1 & 0 & 0 \\
  1 & 0 & 1 & 0 & 1 & 0 \\
  1 & 0 & 0 & 1 & 1 & 0 \\
  1 & 1 & 0 & 0 & 0 & 1 \\
  1 & 0 & 1 & 0 & 0 & 1 \\
  1 & 0 & 0 & 1 & 0 & 0 \\
  1 & 1 & 0 & 0 & 0 & 1 \\
  1 & 0 & 1 & 0 & 0 & 1 \\
  1 & 0 & 0 & 1 & 0 & 0
\end{bmatrix}
\begin{bmatrix}
  \mu \\
  \gamma_1 \\
  \gamma_2 \\
  \gamma_3 \\
  \delta_1 \\
  \delta_2 \\
  \delta_3 \\
  \varepsilon_{11} \\
  \varepsilon_{21} \\
  \varepsilon_{31} \\
  \varepsilon_{12} \\
  \varepsilon_{13} \\
  \varepsilon_{22} \\
  \varepsilon_{32} \\
  \varepsilon_{33}
\end{bmatrix}
\]

(43)

7. The Commutation Matrix

The transformation \(X^c = \odot X^c\) from \(X^c = (x_{ij}e_{jt})\) to \(X^c = (x_{ij}e_{jt})\) is effected by the commutation matrix \(\odot = (e^t_{ij})\). Thus

\[
(x_{ij}e_{jt}) = (e^t_{ij})(x_{ij}e_{jt}).
\]

(44)

It is easy to see that \(\odot\) is an orthonormal matrix with \(\odot' = \odot^{-1}\).

Any pair of covariant indices that are found within a multiple tensor product may be commuted. Thus, for example, there is

\[
\odot(A \otimes B)\odot = (B \otimes A),
\]

(45)

which is expressed in the index notation by writing

\[
(e^k_i)(\{a_{ki}b_{ij}\}e^j_{ki})(e^j_{ti}) = (\{b_{ij}a_{ki}\}e^j_{ik}).
\]

(46)

Here, the four indices may have differing ranges and, therefore, the two commutation matrices of equation (45) may have differing structures and orders. Nevertheless, we have not bothered to make any notational distinction between them. If such a distinction were required, then it might be achieved by attaching to the symbols the indices that are subject to commutations.

Observe also that the braces, which have been applied to the scalar product \(\{a_{ki}b_{ij}\}\) in equation (46), are redundant, since the string contains no repeated indices. Their only effect it to enhance the legibility.

Matrices that perform multiple commutations have been defined by Abadir and Magnus (2005) and by Turkington (2002). Such matrices can be composed
from elementary commutation matrices, such as those defined above. The effects of such transformations is to permutate the sequence of the indices associated with the bases; and, therefore, they can be described in a straightforward manner using the index notation.

**Example.** With \( t \in [1,T] \) and \( j \in [1,M] \), one can write

\[
(47) \quad \odot = (e^j \otimes I_T \otimes e_j) = [I_T \otimes e_1, I_T \otimes e_2, \ldots, I_T \otimes e_M]\]

and

\[
(48) \quad \odot = (e_t \otimes I_M \otimes e^t) = \begin{bmatrix}
I_M \otimes e^1 \\
I_M \otimes e^2 \\
\vdots \\
I_M \otimes e^T
\end{bmatrix}.
\]

The expressions on the RHS of (47) and (48) partly conceal the indices; and they belie the simple nature of the commutation transformation.

**8. Rules for Decomposable Tensor Products**

The following rules govern the decomposable tensors product of matrices, which are commonly described as Kronecker products:

(i) \((A \otimes B)(C \otimes D) = AC \otimes BD\),

(ii) \((A \otimes B)' = A' \otimes B'\),

\[
(49) \quad (iii) \quad A \otimes (B + C) = (A \otimes B) + (A \otimes C),
\]

(iv) \(\lambda(A \otimes B) = \lambda A \otimes B = A \otimes \lambda B\),

(v) \((A \otimes B)^{-1} = (A^{-1} \otimes B^{-1})\).

The Kronecker product is non-commutative, which is to say that \(A \otimes B \neq B \otimes A\). However, observe that

\[
(50) \quad A \otimes B = (A \otimes I)(I \otimes B) = (I \otimes B)(A \otimes I).
\]

This result follows immediately from (49, i).

Such lists of the rules of the Kronecker algebra have been provided in numerous textbooks of econometrics. Examples are the books of Theil (1971) and of Greene (2000) where, in both cases, the results are stated without proofs and without references. One is bound to wonder where the primary sources are to be found. Although the origins of the Kronecker product has been partially uncovered by Henderson *et al.* (1983), its archaeology awaits a full exploration.

To establish the rules of (49) by examples is laborious. However, they can be established readily using the index notation. Consider (49, i), for example, and let

\[
(51) \quad A = (a_{ki}e^i_k), \quad B = (b_{jl}e^j_l), \quad C = (c_{ir}e^r_i), \quad D = (d_{is}e^s_i).
\]
Then,
\[
(A \otimes B)(C \otimes D) = (a_{ki}b_{jl}e_{kl}^{ij})(c_{ir}d_{ts}e_{ir}^{ts})
\]
\[
= (\{a_{ki}c_{ir}\}e_{kr}^i) \otimes (\{b_{jl}d_{ts}\}e_{js}^t)
\]
\[
= AC \otimes BD.
\]

9. Generalised Vectorisation Operations

Turkington (2000 and 2002) has described two matrix operators, called the vec$_n$ and the devec$_m$ operators, which are generalisations, respectively, of the vec operator, which converts a matrix to a long column vector, and of the devec operator, which converts a matrix to a long row vector.

The effect of the vec$_n$ operator is to convert an $m \times np$ partitioned matrix $A = [A_1, A_2, \ldots, A_p]$ to a $pm \times n$ matrix in which the submatrix $A_i$ stands on the shoulders of the submatrix $A_{i+1}$. The devec$_m$ operator is the inverse of the vec$_n$ operator; and its effect would be to convert the vertical stack of matrices into the horizontal array.

To reveal the nature of the vec$_n$ and devec$_m$ operators, let us consider a matrix $A = [a_{ijk}]$, where $i \in [1, \ldots, m]$, $j \in [1, \ldots, n]$ and $k \in [1, \ldots, p]$. If $i$ is a row index and $k, j$ are (lexicographically ordered) column indices, then the matrix, which is of order $m \times np$, has a tensor structure that is revealed by writing
\[
A = (a_{ijk}e_{i}^{kj}).
\]

The effect of the vec$_n$ operator is to convert the index $k$ from a column index to a row index. Thus
\[
\text{vec}_n A = (a_{ijk}e_{i}^{kj})^\tau = (a_{ijk}e_{ki}^{ij}).
\]

Here, the superscripted $\tau$ is an alternative notation for the operator.

The inverse of the vec$_n$ operator is the devec$_m$ operator. Its effect is revealed by writing
\[
\text{devec}_m \{\text{vec}_n(A)\} = (a_{ijk}e_{ki}^{ij})^\tau = (a_{ijk}e_{ki}^{ij}) = A.
\]

Here, the superscripted $\bar{\tau}$ is the alternative notation for the operator.

**Example.** Turkington has provided numerous examples to illustrate the generalised vectorisation operators. Amongst them is the result that
\[
\text{vec}_s[A(E \otimes D)] = (I_q \otimes A)(\text{vec}E \otimes D).
\]

The matrices here are
\[
A = (a_{ijk}e_{i}^{kj}) \quad i \in [1, m], \quad j \in [1, n], \quad k \in [1, p],
\]
\[
D = (d_{jfe}e_{f}^{je}) \quad j \in [1, n], \quad f \in [1, s],
\]
\[
E = (\varepsilon_{kg}e_{k}^{g}) \quad k \in [1, p], \quad g \in [1, q].
\]
D.S.G. POLLOCK

On the LHS of (56), there is

\[
\text{vec}_s[A(E \otimes D)] = \left( (a_{ijk} \epsilon^{kj}_i)(\epsilon_{kg} d_{jf} e^f_{gk_j}) \right)^T \\
= (\{a_{ijk} \epsilon_{kg} d_{jf} \} e^f_{gk_j})^T.
\]

On the RHS of (56), there is

\[
(I_q \otimes A)(\text{vec}E \otimes D) = \left( (a_{ijk} \epsilon^{gkj}_i)(\epsilon_{kg} d_{jf} e^f_{gk_j}) \right) \\
= (\{a_{ijk} \epsilon_{kg} d_{jf} \} e^f_{gk_j}).
\]

Here, the matrix \( I_q = (e^g_q) \), with \( g \in [1, q] \), is embedded within the first parentheses on the RHS without the accompaniment of any scalar terms. (Observe that one might also write \( I_q = (\delta_{g\gamma} e^g_{\gamma}) \), where \( \delta_{g\gamma} \) is Kronecker’s delta.) The equality of (56) follows immediately.

10. The Concept of a Matrix Derivative

Despite the insistence of the majority of the authors who have been cited in the introduction, who have favoured the so-called vectorial definition of the matrix function \( Y = Y(X) \) with respect to its matrix argument \( X \), much use continues to be made of alternative definitions. This circumstance has been noted recently by Magnus (2010), who has demonstrated various deleterious consequences of adopting the alternative non-vectorial definitions.

Amongst the alternative definitions that may be encountered is one that can be specified by writing

\[
\left[ \frac{\partial y_{kl}}{\partial X} \right] = \left( \frac{\partial y_{kl}}{\partial x_{ij}} e^l_{ij} e^k_{ki} \right).
\]

This matrix has the same structure as the product \( Y \otimes X = (y_{kl} x_{ij} e^l_{ij} e^k_{ki}) \), which provides a reasonable recommendation for its use.

A principal reason for using the algebra of Kronecker products and the associated vectorisation operator in multivariate statistical analysis is to allow relationships that are naturally expressed in terms of matrices to be cast into the formats of vector equations. This is to enable the statistical techniques that are appropriate to vector equations to be applied to the matrix equations. To this end, the definition of a matrix derivative should conform to the same algebra as the derivative of a vector function of a vector.

It is commonly agreed that the derivative of the vector function \( y = Ax \) is respect of the vector \( x \) should be the matrix \( \partial y/\partial x = A \), and that of the scalar function \( q = x'Ax \) should be \( \partial q/\partial x = 2x'A \). If analogous definitions are to be applied to matrix equations, then there must be a consistent rule of vectorisation.

Given the pre-existing definition of the Kronecker product of two matrices, which depends upon the lexicographic ordering of the indices, the scope for
alternative methods of vectorisation is strictly limited. The equation $Y = AXB'$ can be vectorised usefully in only two ways, which are represented by equations (35) and (36):

(61) \[ Y^c = (AXB')^c = (B \otimes A)X^c \quad \text{and} \]

(62) \[ Y^\gamma = (AXB')^\gamma = (A \otimes B)X^\gamma, \quad \text{where} \quad Y^\gamma = Y'^c. \]

Then, there are $\partial Y^c/\partial X^c = B \otimes A$ and $\partial Y^\gamma/\partial X^\gamma = A \otimes B$. Both definitions are viable, and both can lead to a fully-fledged theory of matrix differential calculus.

The advantage of (62) over (61) is that the elements of $Y$ are arrayed within the long column vector $Y^\gamma$ according to the lexicographic ordering of their two indices, which is how the elements of $Y$ are liable to be stored within contiguous cells of a computer’s memory.

The elements within the long column vector $Y^c$ follow a reversed lexicographic ordering, which is mildly inconvenient. However, there are also some minor advantages that can be attributed to (61), which represents the conventional method of vectorisation. Thus, the canonical derivative of the matrix function $Y = Y(X)$ with respect to matrix argument $X$ is the vectorial derivative

(63) \[ \frac{\partial Y^c}{\partial X^c} = \left( \frac{\partial y_{ki}}{\partial x_{ij}} c_{jk} \right). \]

Once a multivariate statistical model has been cast in a vectorised format, there remains the task of estimating its parameters. This is commonly approached via the method of maximum-likelihood estimation or via some related method that requires the optimisation of a criterion function. At this stage, it becomes crucial to have a workable theory of matrix differential calculus.

Typically, it is required to differentiate quadratic functions, traces and determinants in respect of matrix arguments. The essential methods of matrix differentiation also comprise product rules and chain rules. The necessary results have been codified in the sources that have been cited in the introduction. Most of these are readily accessible; and it is unnecessary to repeat the results here.

It is in devising an efficient computer code for solving the estimating equations that many of the difficulties of multivariate models can arise. The estimating equations usually entail the Kronecker products of matrices together with commutation transformations and vectorisation operators. In interpreting these, programmers often perform wholly unnecessary and time-wasting manipulations that shift the data between memory cells.

Such operations should be performed wholly within the realms of a pointer arithmetic that serves to provide access to the data objects without the need to shift them from one memory location to another. It is in connection with such a pointer arithmetic that the advantages of the index notation that has been presented in this paper come to the fore.
11. Tensor Products and the Computer

The index notation that has been described in the foregoing sections lends itself readily to an implementation on a computer. The matters arising concern the storage of the data, its recovery and the subsequent computations.

The typical way in which computer memory is allocated to a matrix at run time employs what is commonly described as double indirection. That is to say, a pointer is established that contains the address of an array of pointers, each of which holds the address of the leading element of a matrix row. The memory allocation function ensures that a succession of addresses are linked to that of the leading element of the matrix row, and that these are sufficient to accommodate all of its succeeding elements. A canonical example of such memory allocation is provided by the matrix() function in the book of Numerical Recipes in C of Press et al. (2002).

Such a structure is capable of giving very rapid access to the elements of a matrix. Its disadvantage, from the point of view of tensor algebra, is that it imposes a specific contravariant structure on the tensor product, when it is intended that the form should be mutable. Thus, interest may extend beyond the matrix $X = (x_{ij} e^j_i)$, which is a contravariant tensor product, to its transpose $X^t = (x_{ij} e^j_i)$ and to the vectors $X^c = (x_{ij} e_j^i)$ and $X^r = (x_{ij} e^{ij})$.

To make these forms easily accessible, it is appropriate to store the tensor product as a long vector $X^c$ or $X^r$. It seems preferable to adopt the row vector $X^r = (x_{ij} e^{ij})$ as the stored form, since its elements $x_{ij}$ are arranged according to the natural the lexicographic ordering of the two indices. For the recovery of the data, it is necessary to know the address of the leading element, and to have a record of the range of the indices $i \in [1, m], j \in [1, n]$, together with a recognition that $i$ precedes $j$ in the lexicographic ordering.

To illustrate the recovery of a tensor structure, let the sequence $x[r]; r = 1, \ldots, mn$ comprise the elements $x_{ij}$ arranged according to the lexicographic ordering of the two indices. Then, the following code would serve to print the matrix $X = (x_{ij} e^j_i)$ in $m$ row and $n$ columns:

\begin{equation}
\text{for } i = 1 \text{ to } m \text{ do}
\begin{align*}
&\text{begin} \\
&\text{for } j = 1 \text{ to } n \text{ do} \\
&\quad \text{Write } x[(i-1)n + j] ; \\
&\quad \text{NewLine} ; \\
&\text{end} ;
\end{align*}
\end{equation}

To print the matrix $X^r = (x_{ij} e^{ij})$, we should interchange the two for statements so that the fast inner loop is now in respect of the index $i \in [1, m]$ and the slower outer loop is in respect of the index $j \in [1, n]$.

The formation of multiple tensor products is no more difficult. We might be concerned to form the product $B \otimes A = \{b_{ij} a_{kl} e^{ij}_{kl}\}$ of equation (27), where $B = (b_{ij} e^j_i)$ has $l \in [1, t], j \in [1, n]$, and $A = (a_{ki} e^k_i)$ has $k \in [1, s], i \in [1, m]$. We should store $B^r = (b_{ij} e^{ij})$ and $A^r = (a_{ki} e^{ki})$ and we should calculate $B^r \otimes A^r = \{b_{ij} a_{kl} e^{ij}_{kl}\}$, to be stored in a sequence $x[r]; r = 1, \ldots, mnst$.
according to the lexicographic ordering of the indices \( l_{j}k_{i} \). Then, the following code would serve to print the matrix \( B \otimes A \) of \( st \) rows and \( mn \) columns:

\[
\begin{align*}
&\text{for } l = 1 \text{ to } t \text{ do} \\
&\quad \text{for } k = 1 \text{ to } s \text{ do} \\
&\quad \quad \text{begin} \\
&\quad \quad \quad \text{for } j = 1 \text{ to } n \text{ do} \\
&\quad \quad \quad \quad \text{for } i = 1 \text{ to } m \text{ do} \\
&\quad \quad \quad \quad \quad \text{Write } x[(l-1)msn + (j-1)ms + (k-1)m + i]; \\
&\quad \quad \quad \quad \quad \text{NewLine}; \\
&\quad \quad \quad \text{end;} \\
&\text{end;}
\end{align*}
\]

To understand this code, we may note that the composite index

\[
\begin{align*}
&\quad r = [(l-1)msn + (j-1)ms + (k-1)m + i] \\
&\quad = [(l-1)n + (j-1)s + (k-1)m + i]
\end{align*}
\]

reflects the ordering of the indices \( l_{j}k_{i} \), and that it is in accordance with the sequence in which the matrix elements are stored. We may also note that the alternative ordering of the indices \( l_{k}j_{i} \), which are associated with the nested \texttt{for} statements, corresponds to the order of the indices of \((B \otimes A)^r = \{b_{lj}a_{ki}\}e^{l_{k}j_{i}}\) which, disregarding the breaks occasioned by the \texttt{NewLine} statement, is the order in which we desire to print the elements.

**Example.** A further example may serve to reaffirm the usefulness of the simple tensor algebra and to hint at the problems in computing to which it can be applied. The example concerns the mapping from the graphics memory of a computer to its monitor.

Consider the case of a personal computer of the 1980’s, of which the monitor was intended primarily for a textual display. The textual symbols were formed within character cells that were grids of pixels in eight rows and eight columns. The pixel rows may be indexed by \( k \in [1, h] \) and the pixel columns may be indexed by \( l \in [1, w] \). Then, the generic lighted pixel can be represented by the matrix \( e^{l}_{k} \), which has a unit in the \( k \)th row and the \( l \)th column and zeros elsewhere.

In the monitor, there were 40 character cells per line of text, and there were 25 lines. The character lines may be indexed by \( i \in [1, m] \) and the character columns by \( j \in [1, n] \). The generic character cell may be represented by the matrix \( e^{j}_{i} \), which serves to indicate the location of the cell within the context of the entire screen.

The location of an individual pixel within the screen can be represented by the Kronecker matrix product \( e^{j}_{i} \otimes e^{l}_{k} = e^{jl}_{ik} \); and the contents of the screen may be represented by \( X = (x_{ijkl}e^{jl}_{ik}) \). With 1000 character cells and with 64 bits devoted to each character, the computer’s monochrome graphics memory amounted to 64,000 bits, albeit that a variety of colours were also available, which added an extra dimension to the tensorial structure. The total memory
of the computer was 64 kilobytes. This dimension determined the name of the computer, which was called the Commodore 64—see Jarrett (1984).

Within the computer’s memory, the graphical bits were numbered in a single sequence beginning with the first bit in the first character cell in the top left corner of the screen. The sequence passed from the final bit in the first character cell (number 64 in the lower right corner) to the first bit of the second character of the line. Thereafter, it passed through successive characters of the line. When the end of the line was reached, the sequence passed to the line below.

It follows that the operation that maps the matrix of the pixels of the computer console into the bits of the graphics memory can be represented as follows:

\[(x_{ijkl}c_{ijkl}) \rightarrow (x_{ijkl}c_{ijkl}).\]

The memory location of the generic pixel \(x_{ijkl}\) is indexed by

\[r = (i - 1)whn + (j - 1)wh + (k - 1)w + l.\]

The computer could also be used for plotting scientific graphs. The problem was how to find, within the graphics memory, a pixel that had the screen coordinates \((x, y)\), where \(x \in [1, 320]\) is a number of pixels representing the horizontal location, measured from the left edge of the screen, and \(y \in [1, 200]\) is the vertical location in terms of a pixel count starting from the top of the screen. This is a matter of finding the indices \(i, j, k\) and \(l\) that correspond to the coordinates \(x\) and \(y\). Once the indices have been found, the function \(r(i, j, k, l)\) of (68) could be used to find the memory location.

The two row indices \(i\) and \(k\) are obtained from the vertical coordinate \(y\) using the identity

\[y = h \times (y \text{ div } h) + (y \text{ mod } h)\]
\[= h \times (i - 1) + k,
\]

which gives

\[i = (y \text{ div } h) + 1 \quad \text{and} \quad k = (y \text{ mod } h).
\]

The column indices \(j\) and \(l\) are obtained from the horizontal coordinate \(x\) using the identity

\[x = w \times (x \text{ div } w) + (x \text{ mod } w)\]
\[= w \times (j - 1) + l,
\]

which gives

\[j = (x \text{ div } w) + 1 \quad \text{and} \quad l = (x \text{ mod } w).
\]
TENSOR PRODUCTS and MATRIX DERIVATIVES

References


